

Weight enumerators, intersection enumerators, and Jacobi polynomials

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Abstract

The intersection enumerator and the Jacobi polynomial in an *arbitrary* genus for a binary code are introduced. Adding the weight enumerator into our discussion, we give the explicit relations among them and give some of their basic properties.

1 Introduction

The weight enumerator plays an important role in coding theory. Gleason [6] initiated the applications of the weight enumerator to the invariant theory of the finite groups and Broué-Enguehard [3] constructed modular forms from the weight enumerators. These results were generalized to higher genera [1, 8, 4, 16, 10]. In [15], Ozeki provided the new notion “Jacobi polynomial” of a code. He stated that this comes out of considerations on various invariants of codes [12, 13, 14] and on Jacobi theta-series [5]. In [11], the notion of the intersection enumerator is given for some computations of extremal codes. In the present paper, we discuss these polynomials in an *arbitrary* genus with future applications in mind and some results in [15] are generalized to the case in higher genera.

We shall recall the coding theory (*cf.* [9, 7]). In the present paper we restrict to the binary case. Let $\mathbf{F}_2 = \{0, 1\}$ be the field of two elements and

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\mathbf{F}_2^n the n -dimensional vector space over \mathbf{F}_2 equipped with the usual inner product

$$u \cdot v = u_1v_1 + \cdots + u_nv_n, \quad u = (u_1, \dots, u_n), \quad v = (v_1, \dots, v_n) \in \mathbf{F}_2^n.$$

For $v_1 = (v_{11}, \dots, v_{1n_1}) \in \mathbf{F}_2^{n_1}$, $v_2 = (v_{21}, \dots, v_{2n_2}) \in \mathbf{F}_2^{n_2}$, we put

$$v_1 \oplus v_2 = (v_{11}, \dots, v_{1n_1}, v_{21}, \dots, v_{2n_2}).$$

We introduce an operation \circ on \mathbf{F}_2^n which is given by

$$u \circ v = (u_1v_1, \dots, u_nv_n).$$

This operation satisfies the associative law $(u_1 \circ u_2) \circ u_3 = u_1 \circ (u_2 \circ u_3)$. We denote by $e_{i_1 \dots i_\ell}$ the element of \mathbf{F}_2^g , whose entry is 1 for the i_1, \dots, i_ℓ -part and 0 otherwise. For example, $e_{12} = (1, 1, 0) \in \mathbf{F}_2^3$. Therefore every non-zero element of \mathbf{F}_2^g can be expressed as $e_{i_1 \dots i_\ell}$ for suitable i_1, \dots, i_ℓ . The zero vector in \mathbf{F}_2^g is denoted by e_0 . The weight $wt(u)$ is the number of non-zero coordinates of u . The intersection number $u * v$ in the sense of [15] is $wt(u \circ v)$ in this paper. We denote by $n_a(u_1, \dots, u_g)$ the number of i such that $a = (u_{1i}, \dots, u_{gi})$ for $a \in \mathbf{F}_2^g$.

A linear code of length n is a linear subspace of \mathbf{F}_2^n . We denote by C^\perp the dual code of C :

$$C^\perp = \{u \in \mathbf{F}_2^n \mid u \cdot v = 0, \forall v \in C\}.$$

If $C = C^\perp$, then C is called self-dual. If $wt(u) \equiv 0 \pmod{4}$, $\forall u \in C$, then C is called doubly even. It is known that a self-dual and doubly even code of length n exists if and only if $n \equiv 0 \pmod{8}$.

For a code C_1 (resp. C_2) of length n_1 (resp. n_2), we denote by $C_1 \oplus C_2$ the direct sum of C_1 and C_2 . In other words,

$$C_1 \oplus C_2 = \{u_1 \oplus u_2 \mid u_1 \in C_1, u_2 \in C_2\}.$$

The (homogeneous) weight enumerator of a code C of length n is

$$W_C(x, y) = \sum_{u \in C} x^{n-wt(u)} y^{wt(u)}.$$

This is the first weight enumerator which will be defined in the next section. The inhomogeneous weight enumerator of a code C of length n is¹ $W_C(X) = W_C(x \leftarrow 1, y \leftarrow X)$.

¹The notation $x \leftarrow X$ means to substitute X for x .

2 Definitions of Polynomials of Codes

We start with giving the definitions of the three different kinds of polynomials of codes. It should be emphasized that the notion of the genus attaches to each polynomial.

Definition 1 Let g be a positive integer and C be a code of length n .

(1) The g -th weight enumerator of C is

$$W_C^{(g)}(\{x_a\}_{a \in \mathbf{F}_2^g}) = \sum_{u_1, \dots, u_g \in C} \prod_{a \in \mathbf{F}_2^g} x_a^{n_a(u_1, \dots, u_g)}.$$

(2) The g -th Jacobi polynomial of C with the reference vector $v \in \mathbf{F}_2^n$ is

$$\begin{aligned} & Jac^{(g)}(C, v | \{X_i\}_{1 \leq i \leq g}, \{X_{k_1 \dots k_\ell}\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}}) \\ &= \sum_{u_1, \dots, u_g \in C} \left(\prod_{1 \leq i \leq g} X_i^{wt(u_i)} \right) \left\{ \prod_{2 \leq \ell \leq g+1} \left(\prod_{\substack{1 \leq k_1 < \dots < k_\ell \leq g+1 \\ u_{g+1} = v}} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \right\}. \end{aligned}$$

(3) The g -th intersection enumerator of C is

$$I_C^{(g)}(\{X_{k_1 \dots k_\ell}\}_{\substack{1 \leq \ell \leq g \\ 1 \leq k_1 < \dots < k_\ell \leq g}}) = \sum_{u_1, \dots, u_g \in C} \left\{ \prod_{1 \leq \ell \leq g} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \right\}.$$

When no confusion occurs, we omit the notation of variables in the polynomials and write as $W_C^{(g)}$, $Jac^{(g)}(C, v)$, and $I_C^{(g)}$. It is well known that the g -th weight enumerator is attracted by the number theory (cf. [3, 4, 16]).

Remark 2 (1) $W_C^{(1)}(x_0 \leftarrow 1, x_1 \leftarrow X) = Jac^{(1)}(C, 0 | X_1 \leftarrow X, X_{12}) = I_C^{(1)}(X_1 \leftarrow X)$.

(2) The number of variables in each polynomial is given by

$$W_C^{(g)}: 2^g,$$

$$Jac^{(g)}(C, v): g + \binom{g+1}{2} + \binom{g+1}{3} + \dots + \binom{g+1}{g+1} = 2^{g+1} - 2,$$

$$I_C^{(g)}: \binom{g}{1} + \binom{g}{2} + \dots + \binom{g}{g} = 2^g - 1.$$

At this point, the difference between $Jac^{(g)}(C, v)$ and $I_C^{(g+1)}$ is that $I_C^{(g+1)}$ contains X_{g+1} in its definition, whereas $Jac^{(g)}(C, v)$ does not.

- (3) $I_C^{(g)}(X_{i_1 \dots i_\ell} \leftarrow 1 \text{ for } \ell \geq 3)$ is the intersection enumerator in genus g in the sense of [11].

We shall give the explicit forms for $g = 1, 2$. For $g = 1$, we have

$$Jac^{(1)}(C, v | X_1 \leftarrow X, X_{12} \leftarrow Z) = \sum_{u \in C} X^{wt(u)} Z^{wt(u \circ v)}.$$

This is the Jacobi polynomial of binary codes dealt in [15]. For $g = 2$, we have

$$Jac^{(2)}(C, v) = \sum_{u_1, u_2 \in C} X_1^{wt(u_1)} X_2^{wt(u_2)} X_{12}^{wt(u_1 \circ u_2)} X_{13}^{wt(u_1 \circ u_3)} X_{23}^{wt(u_2 \circ u_3)} X_{123}^{wt(u_1 \circ u_2 \circ u_3)}$$

where $u_3 = v$ in the right-hand side.

3 Basic Results

The Jacobi polynomial $Jac^{(g)}(C, v)$ has the following expansion².

$$Jac^{(g)}(C, v) = \sum_{\substack{\{m_i\} \\ \{r_{k_1 \dots k_\ell}\}}} b(\{m_i\}_{1 \leq i \leq g}, \{r_{k_1 \dots k_\ell}\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}}) \left(\prod_{1 \leq i \leq g} X_i^{m_i} \right) \\ \times \left\{ \prod_{2 \leq \ell \leq g+1} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g+1} X_{k_1 \dots k_\ell}^{r_{k_1 \dots k_\ell}} \right) \right\}$$

where $b(\{m_i\}_{1 \leq i \leq g}, \{r_{k_1 \dots k_\ell}\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}})$ is the number of $(u_1, \dots, u_g) \in C^g$ satisfying

$$\begin{cases} wt(u_i) = m_i & (1 \leq i \leq g), \\ wt(u_{k_1} \circ \dots \circ u_{k_\ell}) = r_{k_1 \dots k_\ell} & (2 \leq \ell \leq g+1, 1 \leq k_1 < \dots < k_\ell \leq g+1). \end{cases}$$

Here we consider as $u_{g+1} = v$. We have a trivial inequality $wt(u_{j_1} \circ \dots \circ u_{j_{\ell+1}}) \leq wt(u_{i_1} \circ \dots \circ u_{i_\ell})$ for $\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_\ell, j_{\ell+1}\}$. From this observation, we have

²If we use the notation $r_i = wt(u_i)$ instead of m_i , some parts of the descriptions below might be simplified, however, we did not take that way because we need to emphasize the distinction between m_* and r_* . Also we respect the usage of the notation in Ozeki's original paper [15].

Proposition 3 *If there exists some pair $\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_\ell, j_{\ell+1}\}$ for some $\ell \geq 1$ such that*

$$m_{i_1} < r_{j_1 j_2} \text{ for } \ell = 1$$

or

$$r_{i_1 \dots i_\ell} < r_{j_1 \dots j_\ell j_{\ell+1}} \text{ for } \ell \geq 2,$$

then we have $b(\{m_i\}, \{r_{i_1 \dots i_\ell}\}) = 0$.

Here we are assuming the convention $u_{g+1} = v$ and $r_{g+1} = wt(u_{g+1})$.

Proposition 4 *Let C be a code of length n which contains all one vector 1 and $Jac^{(g)}(C, v) = \sum b(\{m_i\}, \{r_{k_1 \dots k_\ell}\}) \prod X_i^{m_i} \prod X_{k_1 \dots k_\ell}^{r_{k_1 \dots k_\ell}}$ be the g -th Jacobi polynomial of C with the reference vector v of weight k . Fix j ($1 \leq j \leq g$). Then it holds*

$$b(\{m_i\}, \{r_{k_1 \dots k_\ell}\}) = b(\{m'_i\}, \{r'_{k_1 \dots k_\ell}\})$$

where

$$m'_i = \begin{cases} m_i & \text{if } i \neq j, \\ n - m_j & \text{otherwise,} \end{cases}$$

and

$$r'_{k_1 \dots k_\ell} = \begin{cases} r_{k_1 \dots k_\ell} & \text{if } i \notin \{k_1, \dots, k_\ell\}, \\ r_{k_1 \dots \hat{i} \dots k_\ell} - r_{k_1 \dots i \dots k_\ell} & \text{otherwise,} \end{cases}$$

and \hat{i} means to exclude i .

Proof. The map $(u_1, \dots, u_i, \dots, u_g) \mapsto (u_1, \dots, 1 - u_i, \dots, u_g)$ is a bijection from $\{(u_1, \dots, u_g) \in C^g \mid wt(u_{i_1} \circ \dots \circ u_{i_\ell}) = r_{k_1 \dots k_\ell}, \forall k_1, \dots, k_\ell\}$ to $\{(u_1, \dots, u_g) \in C^g \mid wt(u_{i_1} \circ \dots \circ u_{i_\ell}) = r'_{k_1 \dots k_\ell}, \forall k_1, \dots, k_\ell\}$. This completes the proof of Proposition 4.

Making successive use of this proposition, we get, for $g = 2$,

$$\begin{aligned} & b(\{m_1, m_2\}, \{r_{12}, r_{13}, r_{23}, r_{123}\}) \\ &= b(\{n - m_1, m_2\}, \{m_2 - r_{12}, k - r_{13}, r_{23}, r_{23} - r_{123}\}) \\ &= b(\{m_1, n - m_2\}, \{m_1 - r_{12}, r_{13}, k - r_{23}, r_{13} - r_{123}\}) \\ &= b(\{n - m_1, n - m_2\}, \{n - m_1 - m_2 + r_{12}, k - r_{13}, k - r_{23}, k - r_{13} - r_{23} + r_{123}\}). \end{aligned}$$

Next, we shall consider the reduction of the intersection enumerator to the inhomogeneous weight enumerator.

Proposition 5 (1) $I_C^{(g+1)}(X_{k_1 \dots k_\ell} \leftarrow 1 \text{ for } \ell \geq 2) = \prod_{i=1}^{g+1} W_C(X_i)$.

(2) $I_C^{(g+1)}(\text{variables } \leftarrow 1 \text{ except } X_i) = |C|^g W_C(X_i)$.

Proof. For (1), we have

$$\begin{aligned}
I_C^{(g+1)}(X_{k_1 \dots k_\ell} \leftarrow 1 \text{ for } \ell \geq 2) &= \sum_{u_1, \dots, u_{g+1} \in C^{g+1}} \left(\prod_{1 \leq i \leq g+1} X_i^{wt(u_i)} \right) \left\{ \prod_{2 \leq \ell \leq g+1} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g+1} 1 \right) \right\} \\
&= \sum_{u_1, \dots, u_{g+1} \in C^{g+1}} \left(\prod_{1 \leq i \leq g+1} X_i^{wt(u_i)} \right) \\
&= \prod_{1 \leq i \leq g+1} \left(\sum_{u_i \in C} X_i^{wt(u_i)} \right) \\
&= \prod_{1 \leq i \leq g+1} W_C(X_i).
\end{aligned}$$

The assertion (2) follows from (1) by specializing variables. This completes the proof of Proposition 5.

We would like to provide the relationship between the weight enumerator and the intersection enumerator. In order to do this, we require the following

Lemma 6 *Let u_1, \dots, u_g be elements of \mathbf{F}_2^n . Then the following hold for $1 \leq \ell \leq g$.*

$$\begin{aligned}
(1) \quad wt(u_{i_1} \circ \dots \circ u_{i_\ell}) &= \sum_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq j_1 < \dots < j_k \leq g}} n_{e_{j_1 \dots j_k}}(u_1, \dots, u_g). \\
(2) \quad n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g) &= \sum_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq j_1 < \dots < j_k \leq g}} (-1)^{k-\ell} wt(u_{j_1} \circ \dots \circ u_{j_k}).
\end{aligned}$$

Proof. (1) The right-hand side counts the coordinate positions m such that

$$u_{i_1, m} = u_{i_2, m} = \dots = u_{i_\ell, m} = 1.$$

This is $wt(u_{i_1} \circ \dots \circ u_{i_\ell})$.

(2) The right-hand side is equal to

$$\sum_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq j_1 < \dots < j_k \leq g}} (-1)^{k-\ell} \left\{ \sum_{\substack{\{j_1, \dots, j_k\} \subset \{m_1, \dots, m_s\} \\ 1 \leq m_1 < \dots < m_s \leq g}} n_{e_{m_1 \dots m_s}}(u_1, \dots, u_g) \right\}.$$

We examine each term in this sum. It is easy to see that $n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)$ appears only once. We calculate the coefficient of $n_{e_{m_1 \dots m_s}}(u_1, \dots, u_g)$ where $\{i_1, \dots, i_\ell\} \subsetneq \{m_1, \dots, m_s\}$. The coefficient is

$$\sum_{k=\ell}^s |\{\{j_1, \dots, j_k\} : \{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \subset \{m_1, \dots, m_s\}\}|$$

and this is

$$\begin{aligned} \sum_{k=\ell}^s (-1)^{k-\ell} \binom{s}{k-\ell} &= \sum_{t=0}^{s-\ell} (-1)^t \binom{s-\ell}{t} \\ &= (1 + (-1))^{s-\ell} \\ &= 0. \end{aligned}$$

Thus, the sum under discussion turns out to consist of essentially only one term, that is, $n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)$. This completes the proof of Lemma 6.

Now, we discuss the relations between the weight enumerator and the intersection enumerator.

Theorem 7 *Let C be a code of length n . Then the following hold.*

(1)

$$\begin{aligned} W_C^{(g)} \left(x_{e_0} \leftarrow 1, x_{e_{j_1 \dots j_k}} \leftarrow \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq j_1 < \dots < j_k \leq g}} X_{i_1 \dots i_\ell} \text{ for } 1 \leq k \leq g \right) \\ = I_C^{(g)}(\{X_{i_1 i_2 \dots i_\ell}\}_{\substack{1 \leq \ell \leq g \\ 1 \leq i_1 < i_2 < \dots < i_\ell \leq g}}). \end{aligned}$$

$$(2) \ x_{e_0}^n I_C^{(g)} \left(X_{j_1 \dots j_k} \leftarrow x_{e_0}^{(-1)^k} \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} x_{e_{i_1 \dots i_\ell}}^{(-1)^{k-\ell}} \text{ for } 1 \leq k \leq g \right) = W_C^{(g)}(\{x_a\}_{a \in \mathbb{F}_2^g}).$$

Proof. (1) We have that

$$\begin{aligned} & W_C^{(g)} \left(x_{e_0} \leftarrow 1, x_{e_{j_1 \dots j_k}} \leftarrow \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} X_{i_1 \dots i_\ell} \right) \\ &= \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq j_1 < \dots < j_k \leq g} \left(\prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} X_{i_1 \dots i_\ell} \right)^{e_{j_1 \dots j_k}(u_1, \dots, u_g)} \\ &= \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq j_1 < \dots < j_k \leq g} X_{i_1 \dots i_\ell}^{\sum_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} n_{e_{j_1 \dots j_k}}} \\ &= \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq j_1 < \dots < j_k \leq g} X_{i_1 \dots i_\ell}^{wt(u_{i_1} \circ \dots \circ u_{i_\ell})} \\ &= I_C^{(g)}. \end{aligned}$$

(2) We first observe a trivial relation

$$n_{e_0}(u_1, \dots, u_g) = n - \sum_{1 \leq i_1 < \dots < i_\ell \leq g} n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g).$$

Then we have that

$$\begin{aligned}
W_C^{(g)} &= \sum_{u_1, \dots, u_g \in C} x_{e_0}^{n_{e_0}(u_1, \dots, u_g)} \prod_{1 \leq i_1 < \dots < i_\ell \leq g} x_{e_{i_1 \dots i_\ell}}^{n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)} \\
&= \sum_{u_1, \dots, u_g \in C} x_{e_0}^{n - \sum_{1 \leq i_1 < \dots < i_\ell \leq g} n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)} \prod_{1 \leq i_1 < \dots < i_\ell \leq g} x_{e_{i_1 \dots i_\ell}}^{n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)} \\
&= x_{e_0}^n \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq i_1 < \dots < i_\ell \leq g} \left(\frac{x_{e_{i_1 \dots i_\ell}}}{x_{e_0}} \right)^{n_{e_{i_1 \dots i_\ell}}(u_1, \dots, u_g)} \\
&= x_{e_0}^n \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq i_1 < \dots < i_\ell \leq g} \left(\frac{x_{e_{i_1 \dots i_\ell}}}{x_{e_0}} \right)^{\sum_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq j_1 < \dots < j_k \leq g}} (-1)^{k-\ell} \text{wt}(u_{j_1} \circ \dots \circ u_{j_k})} \\
&= x_{e_0}^n \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq j_1 < \dots < j_k \leq g} \left\{ \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} \left(\frac{x_{e_{i_1 \dots i_\ell}}}{x_{e_0}} \right)^{(-1)^{k-\ell}} \right\}^{\text{wt}(u_{j_1} \circ \dots \circ u_{j_k})}.
\end{aligned}$$

Making use of

$$\begin{aligned}
\prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} \left(\frac{1}{x_{e_0}} \right)^{(-1)^{k-\ell}} &= \left(\frac{1}{x_{e_0}} \right)^{(-1)^{k-1} \binom{k}{1} + (-1)^{k-2} \binom{k}{2} + \dots + (-1)^{k-k} \binom{k}{k}} \\
&= \left(\frac{1}{x_{e_0}} \right)^{-(-1)^k} \\
&= x_{e_0}^{(-1)^k},
\end{aligned}$$

we continue

$$\begin{aligned}
&x_{e_0}^n \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq j_1 < \dots < j_k \leq g} \left(x_{e_0}^{(-1)^k} \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} x_{e_{i_1 \dots i_\ell}}^{(-1)^{k-\ell}} \right)^{\text{wt}(u_{j_1} \circ \dots \circ u_{j_k})} \\
&= x_{e_0}^n I_C^{(g)} \left(X_{j_1 \dots j_k} \leftarrow x_{e_0}^{(-1)^k} \prod_{\substack{\{i_1, \dots, i_\ell\} \subset \{j_1, \dots, j_k\} \\ 1 \leq i_1 < \dots < i_\ell \leq g}} x_{e_{i_1 \dots i_\ell}}^{(-1)^{k-\ell}} \text{ for } 1 \leq k \leq g \right).
\end{aligned}$$

This completes the proof of Theorem 7.

We give some examples.

$$W_C^{(2)}(x_{e_0} \leftarrow 1, x_{e_2} \leftarrow X_2, x_{e_1} \leftarrow X_1, x_{e_{12}} \leftarrow X_1 X_2 X_{12}) = I_C^{(2)}(X_1, X_2, X_{12}).$$

$$\begin{aligned} W_C^{(3)}(x_{e_0} \leftarrow 1, x_{e_1} \leftarrow X_1, x_{e_2} \leftarrow X_2, x_{e_3} \leftarrow X_3, \\ x_{e_{12}} \leftarrow X_1 X_2 X_{12}, x_{e_{13}} \leftarrow X_1 X_3 X_{13}, x_{e_{23}} \leftarrow X_2 X_3 X_{23}, x_{e_{123}} \leftarrow X_1 X_2 X_3 X_{12} X_{13} X_{23} X_{123}) \\ = I_C^{(3)}(X_1, X_2, X_3, X_{12}, X_{13}, X_{23}, X_{123}). \end{aligned}$$

$$x_{e_0}^n I_C^{(2)}(X_1 \leftarrow \frac{x_{e_1}}{x_{e_0}}, X_2 \leftarrow \frac{x_{e_2}}{x_{e_0}}, X_{12} \leftarrow \frac{x_{e_0} x_{e_{12}}}{x_{e_1} x_{e_2}}) = W_C^{(2)}(x_{e_0}, x_{e_1}, x_{e_2}, x_{e_{12}}).$$

$$\begin{aligned} x_{e_0}^n I_C^{(3)}(X_1 \leftarrow \frac{x_{e_1}}{x_{e_0}}, X_2 \leftarrow \frac{x_{e_2}}{x_{e_0}}, X_3 \leftarrow \frac{x_{e_3}}{x_{e_0}}, X_{12} \leftarrow \frac{x_{e_0} x_{e_{12}}}{x_{e_1} x_{e_2}}, X_{13} \leftarrow \frac{x_{e_0} x_{e_{13}}}{x_{e_1} x_{e_3}}, X_{23} \leftarrow \frac{x_{e_0} x_{e_{23}}}{x_{e_2} x_{e_3}}, \\ X_{123} \leftarrow \frac{x_{e_1} x_{e_2} x_{e_3} x_{e_{123}}}{x_{e_0} x_{e_{12}} x_{e_{13}} x_{e_{23}}}) \\ = W_C^{(3)}(x_{e_0}, x_{e_1}, x_{e_2}, x_{e_{12}}, x_{e_3}, x_{e_{13}}, x_{e_{23}}, x_{e_{123}}). \end{aligned}$$

It is known that the Jacobi forms appear in the Fourier-Jacobi expansion of Siegel modular forms (*cf.* [5, 17, 18]). This is one of the motivations for Ozeki for conducting studies on Jacobi polynomials. To get similar results for the g -th Jacobi polynomial, we consider the $(g+1)$ -th intersection enumerator. For a code of length n , it holds

$$\begin{aligned} I_C^{(g+1)} &= \sum_{u_1, \dots, u_{g+1} \in C} X_1^{wt(u_1)} \dots X_g^{wt(u_g)} X_{g+1}^{wt(u_{g+1})} \prod_{2 \leq \ell \leq g+1} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g+1} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \\ &= \sum_{u_{g+1} \in C} \left\{ \sum_{u_1, \dots, u_g \in C} \prod_{2 \leq \ell \leq g+1} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g+1} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \right\} X_{g+1}^{wt(u_{g+1})} \\ &= \sum_{u_{g+1} \in C} Jac^{(g)}(C, u_{g+1}) X_{g+1}^{wt(u_{g+1})} \\ &= \sum_{r=0}^n \left(\sum_{\substack{v \in C \\ wt(v)=r}} Jac^{(g)}(C, v) \right) X_{g+1}^r. \end{aligned}$$

We have thus obtained the following

Theorem 8 *Let C be a code of length n . Assuming that $X_{g+1} = Y$, we have*

$$I_C^{(g+1)} = \sum_{r=0}^n \left(\sum_{\substack{v \in C \\ wt(v)=r}} Jac^{(g)}(C, v) \right) Y^r.$$

In this theorem, the g -th Jacobi polynomial is obtained from the $(g+1)$ -th weight enumerator. If we ask any relation between those of the same genus, we get the following

Proposition 9 *The g -th Jacobi polynomial of a code C with the zero reference vector is the g -th intersection enumerator, that is,*

$$Jac^{(g)}(C, 0) = I_C^{(g)}.$$

Proof. Since we have $wt(u_{k_1} \circ \dots \circ u_{k_\ell}) = 0$ for $k_\ell = g+1$ from our assumption, it follows that

$$\prod_{2 \leq \ell \leq g+1} \left(\prod_{\substack{1 \leq k_1 < \dots < k_\ell \leq g+1 \\ u_{g+1}=v}} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) = \prod_{2 \leq \ell \leq g} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right).$$

Therefore we have

$$\begin{aligned} Jac^{(g)}(C, 0) &= \sum_{u_1, \dots, u_g \in C} \left(\prod_{1 \leq i \leq g} X_i^{wt(u_i)} \right) \left\{ \prod_{2 \leq \ell \leq g+1} \left(\prod_{\substack{1 \leq k_1 < \dots < k_\ell \leq g+1 \\ u_{g+1}=v}} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \right\} \\ &= \sum_{u_1, \dots, u_g \in C} \left(\prod_{1 \leq i \leq g} X_i^{wt(u_i)} \right) \left\{ \prod_{2 \leq \ell \leq g} \left(\prod_{1 \leq k_1 < \dots < k_\ell \leq g} X_{k_1 \dots k_\ell}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \right) \right\} \\ &= I_C^{(g)}. \end{aligned}$$

This completes the proof of Proposition 9.

Proposition 10 *For $i = 1, 2$, let C_i be a code of length n_i and v_i an element of $\mathbf{F}_2^{n_i}$. Then, we have*

$$Jac^{(g)}(C_1, v_1) Jac^{(g)}(C_2, v_2) = Jac^{(g)}(C_1 \oplus C_2, v_1 \oplus v_2).$$

Proof. The proof is straightforward.

We conclude this section with some explicit calculations. Let H be a code of length 8 defined by the generator matrix

$$\begin{pmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}.$$

This is self-dual and doubly even. For simplicity, we set

$$a = X_1, b = X_2, c = X_{12}, d = X_{13}, e = X_{23}, f = X_{123}.$$

We take the following vectors as the reference vectors:

$$v_1 = (0, 0, 0, 0, 0, 0, 0, 0), \quad v_2 = (1, 1, 1, 1, 0, 0, 0, 0), \quad v_3 = (1, 0, 0, 0, 0, 0, 0, 0).$$

The terms are grouped following the calculation made after Proposition 4.

$$\begin{aligned} Jac^{(2)}(H, v_1) &= I_H^{(2)} = \\ &(a^8 b^8 c^8 + a^8 + b^8 + 1) + 14(a^8 b^4 c^4 + b^4) + 14(a^4 b^8 c^4 + a^4) + 14(a^4 b^4 c^4 + a^4 b^4) + 168a^4 b^4 c^2. \end{aligned}$$

$$\begin{aligned} Jac^{(2)}(H, v_2) &= (a^8 b^8 c^8 d^4 e^4 f^4 + b^8 e^4 + a^8 d^4 + 1) + (a^8 b^4 c^4 d^4 e^4 f^4 + b^4 e^4 + a^8 b^4 c^4 d^4 + b^4) \\ &+ 12(a^8 b^4 c^4 d^4 e^2 f^2 + b^4 e^2) + (a^4 b^8 c^4 d^4 e^4 f^4 + a^4 b^8 c^4 e^4 + a^4 d^4 + a^4) \\ &+ 12(a^4 b^8 c^4 d^2 e^4 f^2 + a^4 d^2) + (a^4 b^4 c^4 d^4 e^4 f^4 + a^4 b^4 e^4 + a^4 b^4 d^4 + a^4 b^4 c^4) \\ &+ 12(a^4 b^4 c^4 d^2 e^2 f^2 + a^4 b^4 d^2 e^2) + 12(a^4 b^4 c^2 d^4 e^2 f^2 + a^4 b^4 c^2 e^2) \\ &+ 12(a^4 b^4 c^2 d^2 e^4 f^2 + a^4 b^4 c^2 d^2) + 12(a^4 b^4 c^2 d^2 e^2 f^2 + a^4 b^4 c^2 d^2 e^2) + 96a^4 b^4 c^2 d^2 e^2 f. \end{aligned}$$

$$\begin{aligned} Jac^{(2)}(H, v_3) &= (a^8 b^8 c^8 d e f + b^8 e + a^8 d + 1) + 7(a^8 b^4 c^4 d e f + b^4 e + a^8 b^4 c^4 d + b^4) \\ &+ 7(a^4 b^8 c^4 d e f + a^4 b^8 c^4 e + a^4 d + a^4) + 7(a^4 b^4 c^4 d e f + a^4 b^4 e + a^4 b^4 d + a^4 b^4 c^4) \\ &+ 42(a^4 b^4 c^2 d e f + a^4 b^4 c^2 e + a^4 b^4 c^2 d + a^4 b^4 c^2). \end{aligned}$$

4 MacWilliams Identity for Inhomogeneous Jacobi Polynomial

In this section, we give the MacWilliams identity for the inhomogeneous Jacobi polynomial of a code.

Theorem 11 For a code C of length n , we have that

$$\begin{aligned} & Jac^{(g)}(C^\perp, v \mid \{X_i\}_{1 \leq i \leq g}, \{X_{k_1 \dots k_\ell}\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}}) \\ &= \frac{1}{|C|^g} Z_0^n \left(\frac{Z_{g+1}}{Z_0} \right)^{wt(v)} Jac^{(g)}(C, v \mid \left\{ \frac{Z_i}{Z_0} \right\}_{1 \leq i \leq g}, \{Z_{k_1 \dots k_\ell}\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}}) \end{aligned}$$

where

$$\begin{aligned} Z_0 &= \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = 0}} \left(X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq t \leq g+1 \\ 1 \leq m_1 < \dots < m_t \leq g+1}} X_{m_1 \dots m_t}^{b_{m_1} \dots b_{m_t}} \right), \\ Z_i &= \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = 0}} \left((-1)^{b_i} X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq t \leq g+1 \\ 1 \leq m_1 < \dots < m_t \leq g+1}} X_{m_1 \dots m_t}^{b_{m_1} \dots b_{m_t}} \right) \quad (1 \leq i \leq g), \\ Z_{g+1} &= \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = 1}} \left(X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq t \leq g+1 \\ 1 \leq m_1 < \dots < m_t \leq g+1}} X_{m_1 \dots m_t}^{b_{m_1} \dots b_{m_t}} \right) \end{aligned}$$

and for $\ell \geq 2$

$$\begin{aligned} & Z_{k_1 \dots k_\ell} \\ &= \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} \left(\sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = \delta_{k_{i_j}, g+1}}} (-1)^{b_{k_{i_1}} + \dots + b_{k_{i_{j-1}}} + (1-b_{g+1})b_{k_{i_j}}} \right. \\ & \quad \left. \cdot X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq t \leq g+1 \\ 1 \leq m_1 < \dots < m_t \leq g+1}} X_{m_1 \dots m_t}^{b_{m_1} \dots b_{m_t}} \right)^{(-1)^{\ell-j}}. \end{aligned}$$

Before proceeding to the proof, we give a remark concerning to the definition of $Z_{k_1 \dots k_\ell}$. If $j = 0$ which corresponds to $\emptyset \subset \{k_1, \dots, k_\ell\}$, the product factor is understood as Z_0 . The integer $b_{g+1} = \delta_{k_{i_j}, g+1}$ is a constant and we

need to distinguish the cases $k_{i_j} = g + 1$ and $k_{i_j} < g + 1$. We note

$$(b_{g+1} - 1)k_{i_j} = \begin{cases} k_{i_j} & (k_{i_j} < g + 1), \\ 0 & (k_{i_j} = g + 1). \end{cases}$$

Proof of Theorem 11. We follow the usual method of proving the MacWilliams identity with the function

$$\begin{aligned} \delta_{C^\perp}(w) &= \begin{cases} 1 & w \in C^\perp \\ 0 & w \notin C^\perp \end{cases} \\ &= \frac{1}{|C|} \sum_{u \in C} (-1)^{u \cdot w} \end{aligned}$$

for $w \in \mathbf{F}_2^n$. We have that

$$\begin{aligned} Jac^{(g)}(C^\perp, v) &= \sum_{v_1, \dots, v_g \in C} \left(\prod X_i^{wt(v_i)} \right) \left(\prod X_{k_1 \dots k_\ell}^{wt(v_{k_1} \circ \dots \circ v_{k_\ell})} \right) \\ &= \sum_{w_1, \dots, w_g \in \mathbf{F}_2^n} \left(\prod \delta_{C^\perp}(w_i) X_i^{wt(w_i)} \right) \left(\prod X_{k_1 \dots k_\ell}^{wt(w_{k_1} \circ \dots \circ w_{k_\ell})} \right) \\ &= \sum_{w_1, \dots, w_g \in \mathbf{F}_2^n} \left(\prod \frac{1}{|C|} \sum_{u_i \in C} (-1)^{u_i \cdot w_i} X_i^{wt(w_i)} \right) \left(\prod X_{k_1 \dots k_\ell}^{wt(w_{k_1} \circ \dots \circ w_{k_\ell})} \right) \\ &= \frac{1}{|C|^g} \sum_{\substack{u_1, \dots, u_g \in C \\ w_1, \dots, w_g \in \mathbf{F}_2^n \\ w_{g+1} = v}} (-1)^{(u_{11}w_{11} + \dots + u_{1n}w_{1n}) + \dots + (u_{g1}w_{g1} + \dots + u_{gn}w_{gn})} \\ &\quad \cdot X_1^{w_{11}} \dots X_1^{w_{1n}} \dots X_g^{w_{g1}} \dots X_g^{w_{gn}} \\ &\quad \cdot X_{12}^{w_{11}w_{21}} \dots X_{12}^{w_{1n}w_{2n}} \dots \\ &\quad \cdot X_{k_1 \dots k_\ell}^{w_{k_1 1} \dots w_{k_\ell 1}} \dots X_{k_1 \dots k_\ell}^{w_{k_1 n} \dots w_{k_\ell n}} \\ &\quad \cdot X_{12 \dots g, g+1}^{w_{11}w_{21} \dots w_{g1}w_{g+1,1}} \dots X_{12 \dots g, g+1}^{w_{1n}w_{2n} \dots w_{gn}w_{g+1,n}} \end{aligned}$$

$$\begin{aligned}
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \\
&\prod_{1 \leq i \leq g} \left\{ \sum_{w_{1i}, \dots, w_{gi} \in \mathbf{F}_2} (-1)^{u_{1i}w_{1i} + \dots + u_{gi}w_{gi}} \left(X_1^{w_{1i}} \dots X_g^{w_{gi}} \prod_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} X_{k_1 \dots k_\ell}^{w_{k_1 i} \dots w_{k_\ell i}} \right) \right\} \\
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \\
&\prod_{a \in \mathbf{F}_2^{g+1}} \left\{ \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = a_{g+1}}} (-1)^{a_1 b_1 + \dots + a_g b_g} \left(X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} X_{k_1 \dots k_\ell}^{b_{k_1} \dots b_{k_\ell}} \right) \right\}^{n_a(u_1, \dots, u_g, v)} \\
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \prod_{a \in \mathbf{F}_2^{g+1}} x_a^{n_a(u_1, \dots, u_g, v)}
\end{aligned}$$

where

$$x_a = \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = a_{g+1}}} (-1)^{a_1 b_1 + \dots + a_g b_g} X_1^{b_1} \dots X_g^{b_g} \prod_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} X_{k_1 \dots k_\ell}^{b_{k_1} \dots b_{k_\ell}}.$$

We observe the identity

$$\prod_{a \in \mathbf{F}_2^{g+1}} x_a^{n_a(u_1, \dots, u_g, v)} = x_0^n \prod_{\substack{1 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \left\{ \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} x_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right\}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})}$$

by Lemma 6 (2). Therefore we have

$$\begin{aligned}
Jac(C^\perp, v) &= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \prod_{a \in \mathbf{F}_2^{g+1}} x_a^{n_a(u_1, \dots, u_g, v)} \\
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} x_0^n \prod_{\substack{1 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \left\{ \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} x_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right\}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \\
&= \frac{1}{|C|^g} x_0^n \left(\frac{x_{g+1}}{x_0} \right)^{wt(v)} \\
&\quad \cdot \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq i \leq g} \left(\frac{x_i}{x_0} \right)^{wt(u_i)} \prod_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \left\{ \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} x_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right\}^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})}.
\end{aligned}$$

This is the expected formula stated in Theorem 11 if we put

$$Z_i = x_i \quad (0 \leq i \leq g+1),$$

and for $\ell \geq 2$

$$Z_{k_1 \dots k_\ell} = \prod_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \left\{ \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} x_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right\}.$$

This completes the proof of Theorem 11.

We give the explicit forms for small genera.

For $g = 1$, we have that

$$\begin{aligned}
x_0 &= 1 + X_1, \\
x_1 &= 1 - X_1, \\
x_2 &= 1 + X_1 X_{12}, \\
x_{12} &= 1 - X_1 X_{12}
\end{aligned}$$

and

$$Z_i = x_i, \quad Z_{12} = \frac{x_0 x_{12}}{x_1 x_2}.$$

For $g = 2$, we have that

$$\begin{aligned} x_0 &= 1 + X_1 + X_2 + X_1 X_2 X_{12}, \\ x_1 &= 1 - X_1 + X_2 - X_1 X_2 X_{12}, \\ x_2 &= 1 + X_1 - X_2 - X_1 X_2 X_{12}, \\ x_3 &= 1 + X_1 X_{13} + X_2 X_{23} + X_1 X_2 X_{12} X_{13} X_{23} X_{123}, \\ x_{12} &= 1 - X_1 - X_2 + X_1 X_2 X_{12}, \\ x_{13} &= 1 - X_1 X_{13} + X_2 X_{23} - X_1 X_2 X_{12} X_{13} X_{23} X_{123}, \\ x_{23} &= 1 + X_1 X_{13} - X_2 X_{23} - X_1 X_2 X_{12} X_{13} X_{23} X_{123}, \\ x_{123} &= 1 - X_1 X_{13} - X_2 X_{23} + X_1 X_2 X_{12} X_{13} X_{23} X_{123} \end{aligned}$$

and

$$Z_i = x_i, \quad Z_{12} = \frac{x_0 x_{12}}{x_1 x_2}, \quad Z_{13} = \frac{x_0 x_{13}}{x_1 x_3}, \quad Z_{23} = \frac{x_0 x_{23}}{x_2 x_3}, \quad Z_{123} = \frac{x_1 x_2 x_3 x_{123}}{x_0 x_{12} x_{13} x_{23}}.$$

5 MacWilliams Identity for Homogeneous Jacobi Polynomial

In this section, we give the MacWilliams identity for the homogeneous Jacobi polynomial of a code.

In order to homogenize the Jacobi polynomial, we introduce the new variables

$$X_{k_1 \dots k_\ell} = \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} y_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}}.$$

Then the homogeneous Jacobi polynomial $\mathfrak{Jac}(C, v; \{y_a\}_{a \in \mathbb{F}_2^{g+1}})$ is defined as

follows.

$$\begin{aligned}
& \mathfrak{Jac}^{(g)}(C, v; \{y_a\}_{a \in \mathbf{F}_2^{g+1}}) \\
&= y_0^n \left(\frac{y_{g+1}}{y_0} \right)^{wt(v)} \text{Jac}^{(g)}(C, v; \left\{ \frac{y_i}{y_0} \right\}_{1 \leq i \leq g}), \left\{ \prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} y_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right\}_{\substack{2 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \right) \\
&= \sum_{u_1, \dots, u_g \in C} y_0^n \prod_{\substack{1 \leq \ell \leq g+1 \\ 1 \leq k_1 < \dots < k_\ell \leq g+1}} \left(\prod_{\substack{0 \leq j \leq \ell \\ 1 \leq i_1 < \dots < i_j \leq \ell \\ \{k_{i_1}, \dots, k_{i_j}\} \subset \{k_1, \dots, k_\ell\}}} y_{k_{i_1} \dots k_{i_j}}^{(-1)^{\ell-j}} \right)^{wt(u_{k_1} \circ \dots \circ u_{k_\ell})} \\
&= \sum_{u_1, \dots, u_g \in C} \prod_{a \in \mathbf{F}_2^{g+1}} y_a^{n_a(u_1 \dots u_g v)}
\end{aligned}$$

This is a homogeneous polynomial of total degree n . Next we give the MacWilliams identity for $\mathfrak{Jac}(C, v; \{y_a\})$.

Theorem 12 *For a code C of length n , we have*

$$\mathfrak{Jac}(C^\perp, v, g; \{y_a\}_{a \in \mathbf{F}_2^{g+1}}) = \frac{1}{|C|^g} \mathfrak{Jac}(C, v; \left\{ \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = a_{g+1}}} (-1)^{a_1 b_1 + \dots + a_g b_g} y_{(b_1 \dots b_g b_{g+1})} \right\}_{a \in \mathbf{F}_2^{g+1}}).$$

Proof. We have

$$\begin{aligned}
\mathfrak{Jac}(C^\perp, v; \{y_a\}) &= \sum_{u_1, \dots, u_g \in C^\perp} \prod_{a \in \mathbf{F}_2^{g+1}} y_a^{n_a(u_1, \dots, u_g, v)} \\
&= \sum_{w_1, \dots, w_g \in \mathbf{F}_2^n} \delta_{C^\perp}(w_1) \dots \delta_{C^\perp}(w_g) \prod_a y_a^{n_a(w_1, \dots, w_g, v)} \\
&= \frac{1}{|C|^g} \sum_{\substack{u_1, \dots, u_g \in C \\ w_1, \dots, w_g \in \mathbf{F}_2^n}} (-1)^{u_1 \cdot w_1 + \dots + u_g \cdot w_g} \prod_a y_a^{n_a(w_1, \dots, w_g, v)} \\
&= \frac{1}{|C|^g} \sum_{\substack{u_1, \dots, u_g \in C \\ w_1, \dots, w_g \in \mathbf{F}_2^n}} (-1)^{(u_{11}w_{11} + \dots + u_{1n}w_{1n}) + \dots + (u_{g1}w_{g1} + \dots + u_{gn}w_{gn})} \\
&\quad \times y(w_{11} \dots w_{g1} v_1) \dots y(w_{1n} \dots w_{gn} v_n) \\
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \prod_{1 \leq i \leq n} \left\{ \sum_{w_{1i}, \dots, w_{gi} \in \mathbf{F}_2} (-1)^{u_{1i}w_{1i} + \dots + u_{gi}w_{gi}} y(w_{1i} \dots w_{gi} v_i) \right\} \\
&= \frac{1}{|C|^g} \sum_{u_1, \dots, u_g \in C} \prod_{a \in \mathbf{F}_2^{g+1}} \left\{ \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = a_{g+1}}} (-1)^{a_1 b_1 + \dots + a_g b_g} y(b_1 \dots b_g b_{g+1}) \right\}^{n_a(u_1, \dots, u_g, v)} \\
&= \frac{1}{|C|^g} \mathfrak{Jac}(C, v; \left\{ \sum_{\substack{b_1, \dots, b_g \in \mathbf{F}_2 \\ b_{g+1} = a_{g+1}}} (-1)^{a_1 b_1 + \dots + a_g b_g} y(b_1 \dots b_g b_{g+1}) \right\}_{a \in \mathbf{F}_2^{g+1}}).
\end{aligned}$$

This completes the proof of the MacWilliams identity for the homogeneous Jacobi polynomial.

We will discuss the applications of the MacWilliams identity in the subsequent papers.

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